

Mixing of Quantum Walks on Generalized Hypercubes

Ana Best
McGill University

Markus Kliegl
Swarthmore College

Shawn Mead-Gluchacki
SUNY Potsdam

Christino Tamon*
Clarkson University

August 18, 2008

Abstract

We study continuous-time quantum walks on graphs which generalize the hypercube. The only known family of graphs whose quantum walk instantaneously mixes to uniform is the Hamming graphs with small arities. We show that quantum uniform mixing on the hypercube is robust under the addition of perfect matchings but not much else. Our specific results include:

- The graph obtained by augmenting the hypercube with an additive matching $x \mapsto x \oplus \eta$ is instantaneous uniform mixing whenever $|\eta|$ is even, but with a slower mixing time. This strictly includes Moore-Russell's result on the hypercube [11].
- The class of Hamming graphs $H(n, q)$ is not uniform mixing if and only if $q \geq 5$. This is a tight characterization of quantum uniform mixing on Hamming graphs; previously, only the status of $H(n, q)$ with $q < 5$ was known.
- The bunkbed graph $\mathcal{B}_n(A_f)$ whose adjacency matrix is $I \otimes Q_n + X \otimes A_f$, where A_f is a \mathbb{Z}_2^n -circulant matrix defined by a Boolean function f , is not uniform mixing if the Fourier transform of f has support of size smaller than 2^{n-1} . This explains why the hypercube is uniform mixing and why the join of two hypercubes is not.

Our work exploits the rich spectral structure of the generalized hypercubes and relies heavily on Fourier analysis of group-circulants.

Keywords: Quantum walk, continuous-time, mixing, group-circulant, hypercube.

1 Introduction

Quantum walk on graphs is an important area of study in quantum information and computation for at least two reasons. The first is algorithmic; as a natural generalization of classical random walks, quantum walks offer an alternative paradigm to develop new quantum algorithms (see [4, 8]). The second reason is physical; quantum walks provide a potentially simpler method for implementing quantum computers. Although both arguments are arguably still being debated and investigated, research into the intrinsic properties of quantum walk continues to be an important step towards these algorithmic and physical goals.

*Contact author: tino@clarkson.edu

An excellent survey on quantum walks is given by Kendon [10]. In this paper, we focus on the continuous-time quantum walk on finite graphs and study its mixing properties.

In their breakthrough paper [11], Moore and Russell showed that the continuous-time quantum walk on the hypercube is instantaneous uniform mixing; more importantly, it mixes faster than the classical walk. Most well-known graphs, however, do not exhibit quantum uniform mixing, as later results showed. These include graphs such as the complete graphs, the Cayley graphs of the symmetric group, and most even-length cycles (see [1, 9, 2]). In contrast, classical random walks are uniform mixing on most well-behaved graphs. This suggests that, unlike the classical case, uniform mixing is a rare phenomenon in quantum walks, and that the hypercube is an anomaly. We investigate if this quantum uniform mixing phenomenon still exists in some natural generalizations of the hypercube.

First, we consider the standard hypercube augmented with perfect matchings defined by additive shifts $x \mapsto x \oplus \eta$, for $\eta \in \mathbb{Z}_2^n$. How destructive are these added matchings to uniform mixing on the hypercube? Surprisingly, we found that the resulting hypercube Q_n^η is *still* instantaneous uniform mixing whenever $|\eta|$ is even, and, more interestingly, that it has a *slower* mixing time than the standard hypercube. Since the augmented matchings lower the diameter of the hypercube, this is yet another example of a counter-intuitive phenomenon in quantum walks. This shows that quantum uniform mixing on the hypercube is robust under additive matchings. For $|\eta|$ odd, we discover that the quantum walk is uniform mixing provided it starts in the superposition $\frac{1}{\sqrt{2}}(|0_n\rangle + |\eta\rangle)$.

Second, we study the class of Hamming graphs $H(q, n)$ which are n -dimensional q -ary hypercubes (see Biggs [3]). Prior to this work, it was known that $H(q, n)$ is quantum uniform mixing if $q \in \{2, 3, 4\}$ (see [5]). We close this gap by showing that $H(q, n)$ is *not* quantum uniform mixing if $q \geq 5$, for any n . This gives a tight characterization of quantum uniform mixing on Hamming graphs.

Third, we generalize the hypercube using its recursive construction by Cartesian products. The $(n + 1)$ -dimensional hypercube Q_{n+1} is built by combining two n -dimensional hypercubes Q_n by connecting their corresponding vertices. More formally, $Q_{n+1} = Q_n \oplus K_2$, where K_2 is the complete graph on two vertices. The adjacency matrix of this Cartesian product is given by $A_{Q_{n+1}} = I \otimes A_{Q_n} + X \otimes I$. So, the connection between the two copies of Q_n is specified by the (second) identity matrix I . We explore the effect on quantum uniform mixing when this connection is varied.

To this end, we consider the bunkbed graph $\mathcal{B}_n(A_f)$ whose adjacency matrix is $I \otimes A_{Q_n} + X \otimes A_f$, where A_f is a \mathbb{Z}_2^n -circulant defined by a Boolean function f over \mathbb{Z}_2^n (see Diaconis [6]). Much like a standard circulant matrix, the function f defines the first row of the matrix A_f and the group operation of \mathbb{Z}_2^n determines the rest of the rows of A_f . So, A_f defines (and generalizes) the connection between the two hypercubes Q_n . For example, the standard hypercube $Q_n = \mathcal{B}_{n-1}(I)$ has $f(x) = \delta_{x, 0_n}$ and the hypercube with an additive matching $Q_n^\eta = \mathcal{B}_{n-1}(A_f)$ has $f(x) = 1$ whenever $x = 0_n$ or $x = \tilde{\eta}$, where $\eta = 1 \cdot \tilde{\eta}$ without loss of generality. Our main result is that $\mathcal{B}_n(A_f)$ is not uniform mixing whenever the Fourier transform of f has small support, that is, $|\text{supp}(\hat{f})| < 2^{n-1}$. An immediate corollary shows that the graph-theoretic join of two hypercubes $Q_n + Q_n$ is not uniform mixing. Unfortunately, the small Fourier support size is not a necessary condition since $\mathcal{B}_n(Q_n)$ is not uniform mixing even though $|\text{supp}(\hat{f})| \geq 2^{n-1}$.

Our work exploits the rich spectral structure of the generalized hypercubes and relies

heavily on Fourier analysis of group-circulants. A more complete treatment of the latter may be found in Diaconis [6].

We summarize the known status of quantum uniform mixing on graphs along with our contributions in Figure 1.

2 Preliminaries

For a logical statement \mathcal{S} , the Iversonian $\llbracket \mathcal{S} \rrbracket$ is 1 if \mathcal{S} is true, and 0 otherwise. Let \mathbb{Z}_m denote the additive group of integers $\{0, \dots, m-1\}$ modulo m . For $a, b \in \mathbb{Z}_2^n$, let $a \oplus b$ denote the bit-wise exclusive OR of a and b , let $a \cdot b = \sum_{k=1}^n a_k b_k \pmod{2}$ denote the inner product modulo 2, and let the Hamming weight $|a|$ be the number of ones in a . We let $e_j \in \mathbb{Z}_2^n$ denote the unit vector that is 1 in position j and zero elsewhere. We use I and J to denote the identity and all-one matrices, respectively; we use X to denote the Pauli- σ_X matrix.

The graphs $G = (V, E)$ we study are finite, simple, undirected, and connected. The adjacency matrix A_G of a graph G is defined as $A_G[u, v] = \llbracket (u, v) \in E \rrbracket$. In most cases, we also require G to be vertex-transitive, that is, for any $a, b \in V$, there is an automorphism $\pi \in \text{Aut}(G)$ with $\pi(a) = b$. The Cartesian product $G \oplus H$ of graphs G and H is a graph whose adjacency matrix is $I \otimes A_G + A_H \otimes I$. Let K_n denote the complete graph on n vertices. Then, the binary n -dimensional hypercube Q_n may be defined recursively as $Q_n = Q_{n-1} \oplus K_2$, for $n \geq 2$, and $Q_1 = K_2$. For more background on algebraic graph theory, we refer the reader to Biggs [3].

Next, we describe group-theoretic circulant graphs and Fourier analysis on \mathbb{Z}_2^n [6]. Let \mathcal{G} be a finite group of order m and let $f : \mathcal{G} \rightarrow \mathbb{C}$ be a class function over \mathcal{G} (that is, f is constant on the conjugacy classes of \mathcal{G}). Then, the $m \times m$ matrix defined by $A_f^{\mathcal{G}}[s, t] = f(ts^{-1})$ is called a \mathcal{G} -circulant matrix defined by f . Moreover, $A_f^{\mathcal{G}}$ defines a \mathcal{G} -circulant graph if f is a $\{0, 1\}$ -valued function that satisfies $f(e) = 0$, where e is the identity element, and $f(a^{-1}) = f(a)$, for all $a \in \mathcal{G}$. Here, the correspondence with Cayley graphs is recovered by letting the generator set be $\{a : f(a) = 1\}$. In this paper, we focus on the Abelian group $\mathcal{G} = \mathbb{Z}_2^n$.

Figure 1: Instantaneous Uniform Mixing on Various Graphs.

Family of Graphs	Mixing	Reference
Hamming graph $H(n, q)$, $q \in \{2, 3, 4\}$	Yes	Moore-Russell [11], Carlson <i>et al.</i> [5]
Complete multipartite graph	No	Ahmadi <i>et al.</i> [1]
Symmetric group	No	Gerhardt-Watrous [9]
Cycles	No ^a	Adamczak <i>et al.</i> [2]
Hamming graph $H(n, q)$, $q \geq 5$	No	this work
$\mathcal{B}_n(A_f)$, $ \text{supp}(f) \in \{1, 2\}$	Yes ^b	Moore-Russell [11], this work
$\mathcal{B}_n(A_f)$, $ \text{supp}(\hat{f}) < 2^{n-1}$	No	this work

^aThis was proved for a subclass of even-length cycles

^bFor $|\text{supp}(f)| = 2$, we require $\text{supp}(f) = \{0_n, a\}$ with $|a|$ odd.

$$\begin{bmatrix}
a & b & c & d & e & f & g & h \\
b & a & d & c & f & e & h & g \\
c & d & a & b & g & h & e & f \\
d & c & b & a & h & g & f & e \\
e & f & g & h & a & b & c & d \\
f & e & h & g & b & a & d & c \\
g & h & e & f & c & d & a & b \\
h & g & f & e & d & c & b & a
\end{bmatrix} \tag{1}$$

Figure 2: Example of a \mathbb{Z}_2^n -circulant matrix: $n = 3$; the first row completely determines the matrix through the group operation of \mathbb{Z}_2^n . The hypercube requires $b = c = e = 1$ and $a = d = f = g = h = 0$.

For some basic facts of Fourier analysis over \mathbb{Z}_2^n , let $f, g : \mathbb{Z}_2^n \rightarrow \mathbb{C}$ be arbitrary functions. The inner product of f, g is defined as $\langle f, g \rangle = \sum_x f(x)g^*(x)$. The group characters of \mathbb{Z}_2^n are given by $\chi_a(x) = (-1)^{a \cdot x}$, for $a, x \in \mathbb{Z}_2^n$, and they satisfy $\langle \chi_a, \chi_b \rangle = 2^n \mathbb{I}[a = b]$. With this, we can define the Fourier transform of f at a as

$$\hat{f}(a) = \langle f, \chi_a \rangle = \sum_x f(x)\chi_a(x) \tag{2}$$

while the inverse Fourier transform is given by

$$f(x) = 2^{-n} \sum_a \hat{f}(a)\chi_a(x). \tag{3}$$

The support of f is $\text{supp}(f) = \{x : f(x) \neq 0\}$; similarly, the support of \hat{f} is $\text{supp}(\hat{f}) = \{a : \hat{f}(a) \neq 0\}$. The Convolution Theorem states that

$$\widehat{fg}(a) = \frac{1}{2^n} \sum_b \hat{f}(b)\hat{g}(a \oplus b). \tag{4}$$

If $P : \mathbb{Z}_2^n \rightarrow [0, 1]$ is a probability distribution, then $\hat{P}(0_n) = 1$. Moreover, P is the uniform distribution if and only if $\hat{P}(a) = 0$ for all $a \neq 0_n$.

If $G = (V, E)$ is a graph with adjacency matrix A , let $|\psi(t)\rangle \in \mathbb{C}^{|V|}$ be a time-dependent amplitude vector over V . Then, the continuous-time quantum walk on G is defined using Schrödinger's equation as

$$|\psi(t)\rangle = e^{-itA}|\psi(0)\rangle, \tag{5}$$

where $|\psi(0)\rangle$ is the initial amplitude vector (see [7]). The *instantaneous* probability of vertex v at time t is $p_v(t) = |\langle v|\psi(t)\rangle|^2$. We say G is *instantaneous uniform mixing* if there is a time t^* such that the quantum walk on G satisfies $|\langle v|\psi(t^*)\rangle|^2 = 1/|V|$, for all $v \in V$.

2.1 Fourier Analysis of the Hypercube

We briefly review the Fourier analysis for the hypercube Q_n [11]. Since the adjacency matrix A of Q_n is a \mathbb{Z}_2^n -circulant, its eigenvectors are the characters $|\chi_a\rangle$ (expressed in ket notation) where $\langle x|\chi_a\rangle = \chi_a(x) = (-1)^{a \cdot x}$, for $a, x \in \mathbb{Z}_2^n$. The first row of A is defined by a Boolean function f where $f(x) = \llbracket |x| = 1 \rrbracket$. Thus, the eigenvalues of A are given by $\lambda_a = \sum_x f(x) \chi_a(x) = \sum_{j=1}^n \chi_a(e_j) = n - 2|a|$. This follows from the theory of group circulants (see Diaconis [6]), but can also be verified directly. Since the hypercube is vertex-transitive, we may assume that the start vertex is $|0\rangle$ (which corresponds to vertex 0_n); moreover, $|0\rangle = 2^{-n} \sum_a |\chi_a\rangle$.

The quantum walk on Q_n starting at $|0\rangle$ is given by $|\psi(t)\rangle = e^{-itA}|0\rangle = 2^{-n} \sum_a e^{-it\lambda_a} |\chi_a\rangle$. Viewing time t as being *fixed*, we view the amplitude vector $|\psi(t)\rangle$ as a function of $a \in \mathbb{Z}_2^n$ and redefine $\psi_t(a) = \langle a|\psi(t)\rangle$. By Fourier inversion, we see that $\widehat{\psi}_t(a) = e^{-it\lambda_a}$. Since $P_t(a) = |\langle a|\psi(t)\rangle|^2 = \psi_t^*(a) \psi_t(a)$, using the Convolution Theorem, we have $\widehat{P}_t(a) = 2^{-n} \sum_b \widehat{\psi}_t(b) \widehat{\psi}_t^*(a \oplus b)$. Since $\widehat{\psi}_t^* = \widehat{\psi}_t^\star$, we obtain

$$\widehat{P}_t(a) = \frac{1}{2^n} \sum_b \exp(-it(\lambda_b - \lambda_{a \oplus b})) \quad (6)$$

To show that P_t is uniform, it suffices to show $\widehat{P}_t(a) = \llbracket a = 0_n \rrbracket$; which was proved by Moore and Russell [11].

3 Hypercube with Additive Matchings

In the standard hypercube Q_n , whose vertices are the elements of \mathbb{Z}_2^n , two vertices a, b are adjacent if $a \oplus b = e_j$, for some j ; that is, a and b differ in exactly one coordinate. For $\eta \in \mathbb{Z}_2^n$, we define the $(n+1)$ -regular graph Q_n^η to be the graph obtained from Q_n by adding the matching $(a, a \oplus \eta)$, for all $a \in \mathbb{Z}_2^n$.

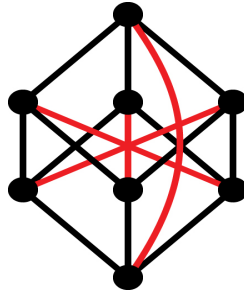


Figure 3: Example of Q_n^η hypercube: $n = 3$ and $\eta = 111$.

Theorem 1 *For $n \geq 2$, a continuous-time quantum walk on Q_n^η is instantaneous uniform mixing if and only if $|\eta|$ is even.*

Proof Let A be the scaled adjacency matrix of Q_n^η , where $A_{a,b} = 1/(n+1)$ if $a \oplus b \in \{e_1, \dots, e_n, \eta\}$, and $A_{a,b} = 0$ otherwise. Let $A_0 = \{a \mid a \cdot \eta = 0\}$ and $A_1 = \{a \mid a \cdot \eta = 1\}$. Note $a \oplus b \in A_0$ if and only if $a, b \in A_0$ or $a, b \in A_1$.

By the group-circulant structure of A , its eigenvalues are given by

$$\lambda_a = \frac{1}{n+1} (n - 2|a| + (-1)^{a \cdot \eta}) = \left(1 - \frac{2k_\eta(a)}{n+1}\right) \quad (7)$$

where $k_\eta(a) = |a| + \llbracket a \in A_1 \rrbracket$. We note that k_η satisfies

$$k_\eta(a \oplus b) - k_\eta(b) = |a \oplus b| - |b| + \llbracket a \in A_1 \rrbracket (-1)^{b \cdot \eta} \quad (8)$$

Since the quantum walk on Q_n^η is given by $|\psi(t)\rangle = \frac{1}{2^n} \sum_a e^{-it\lambda_a} |\chi_a\rangle$, we see that

$$\widehat{\psi}_t(a) = \exp \left[-it \left(1 - \frac{2k_\eta(a)}{n+1} \right) \right]. \quad (9)$$

A similar analysis to Equation (6) yields

$$\widehat{P}_t(a) = \frac{1}{2^n} \sum_b \exp \left[-\frac{2it}{n+1} (k_\eta(a \oplus b) - k_\eta(b)) \right]. \quad (10)$$

By Equation (8), we obtain

$$\widehat{P}_t(a) = \begin{cases} 1 & \text{if } a = 0_n \\ 2^{-n} \sum_b \exp \left[-\frac{2it}{n+1} (|a \oplus b| - |b| + \llbracket a \in A_1 \rrbracket (-1)^{b \cdot \eta}) \right] & \text{otherwise} \end{cases} \quad (11)$$

For $a \in A_0$, analysis similar to that in [11] shows that $\widehat{P}_t(a)$ is periodic in t with period $(n+1)\pi$ and that, up to periodicity, $\widehat{P}_t(a) = 0$ only for $t_1^* = (n+1)\frac{\pi}{4}$ and $t_2^* = (n+1)\frac{3\pi}{4}$. This shows that the only possible times at which Q_n^η could possibly be uniform mixing are t_1^* and t_2^* . We will show for $t = t_1^*$ that we indeed have $\widehat{P}_t(a) = 0$ for all $a \in A_1$ if and only if $|\eta|$ is even. The proof for $t = t_2^*$ works in exactly the same way.

So assume that $a \in A_1$. Define

$$\rho_a(b) = |a \oplus b| - |b| + (-1)^{b \cdot \eta}. \quad (12)$$

Observe that $\rho_a(a \oplus b) = -\rho_a(b)$. Using this symmetry,

$$\widehat{P}_t(a) = \frac{1}{2^n} \sum_b \frac{1}{2} \left[\exp \left(-\frac{2it}{n+1} \rho_a(b) \right) + \exp \left(-\frac{2it}{n+1} \rho_a(a \oplus b) \right) \right] \quad (13)$$

$$= \frac{1}{2^n} \sum_b \cos \left(\frac{2t}{n+1} \rho_a(b) \right) \quad (14)$$

At $t = t_1^*$, we have

$$\widehat{P}_{t_1^*}(a) = \sum_b \cos \left(\frac{\pi}{2} \rho_a(b) \right) \quad (15)$$

Let m be the number of overlaps of 1's between a and b . Then

$$|a \oplus b| - |b| = |a| - 2m. \quad (16)$$

Note that m is even if $a \cdot b = 0$, and m is odd if $a \cdot b = 1$. Using these facts, we see that

$$\sum_b \cos\left(\frac{\pi}{2} p_a(b)\right) = \sum_b \cos\left(\frac{\pi}{2} (|a| - 2m + (-1)^{b \cdot \eta})\right) \quad (17)$$

$$= \sum_b \cos\left(\frac{\pi}{2} (|a| + (-1)^{b \cdot \eta})\right) \cos(m\pi) \quad (18)$$

$$= \sum_b (-1)^{a \cdot b} \cos\left(\frac{\pi}{2} (|a| + (-1)^{b \cdot \eta})\right) \quad (19)$$

If $|a|$ is even, the last expression is 0 and we are done. Therefore assume for the rest of the argument that $|a|$ is odd. Then,

$$\sum_b \cos\left(\frac{\pi}{2} p_a(b)\right) = \sum_b (-1)^{a \cdot b} \cos\left(\pi \frac{(|a| + 1)}{2} + \pi \frac{((-1)^{b \cdot \eta} - 1)}{2}\right) \quad (20)$$

$$= \sum_b (-1)^{a \cdot b} (-1)^{\frac{|a|+1}{2}} (-1)^{b \cdot \eta} \quad (21)$$

$$= (-1)^{\frac{|a|+1}{2}} \sum_b \chi_b(a \oplus \eta) \quad (22)$$

$$= (-1)^{\frac{|a|+1}{2}} \llbracket a = \eta \rrbracket 2^n \quad (23)$$

Finally, notice that, since we are considering $a \in A_1$ such that $|a|$ is odd only, the case $a = \eta$ will occur if and only if $\eta \in A_1$ and $|\eta|$ is odd. Since the condition $\eta \in A_1$ is equivalent to $|\eta|$ being odd, this proves the claim. \square

Remark 1 This includes the Moore-Russell result in the following sense. For $\eta = 0_n$, the Theorem shows that Q_n^η is uniform mixing. If $\eta = 0_n$, however, all we are doing is adding self-loops to the standard hypercube, which just amount to scaling time by the factor $n/(n+1)$. Hence this shows that the standard hypercube is uniform mixing.

Remark 2 The case when $|\eta|$ is odd is interesting as well. Namely, it is near uniform mixing in the following sense. If $|\eta|$ is odd, taking the inverse Fourier transform of \hat{P}_t at time t_1^* or t_2^* yields

$$P_t(a) = \begin{cases} \llbracket a \in A_0 \rrbracket 2^{-n+1} & \text{if } |\eta| \equiv 3 \pmod{4} \\ \llbracket a \in A_1 \rrbracket 2^{-n+1} & \text{if } |\eta| \equiv 1 \pmod{4} \end{cases} \quad (24)$$

That is, the probability distribution at these times is uniform on half of the vertices and zero on the other half. Moreover, it can be shown that if the quantum walk is started in the superposition $|\psi(0)\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |\eta\rangle)$, then P_t is uniform at times t_1^* and t_2^* , just as in the case when $|\eta|$ is even.

4 Hamming Graphs

As defined in Biggs [3], the *Hamming* graph $H(n, q)$ has vertex set $V = \{1, \dots, q\}^n$, and two vertices are adjacent if they differ in exactly one coordinate. As such, Hamming graphs are a natural generalization of the hypercubes Q_n (which are simply $H(n, 2)$).

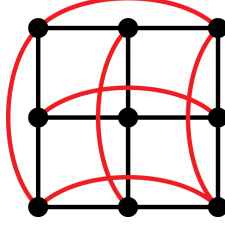


Figure 4: Example of a Hamming graph: $H(2, 3)$.

We prove that most of these Hamming graphs are not instantaneous uniform mixing. Before we prove this result, however, we need a preliminary result about quantum walks on the Cartesian product of graphs, which in turn relies on the following Fact.

Fact 2 *A quantum walk on $G \oplus H$ starting on vertex (g, h) satisfies*

$$|\psi_{G \oplus H}(t)\rangle = |\psi_G(t)\rangle \otimes |\psi_H(t)\rangle, \quad (25)$$

where ψ_G and ψ_H are quantum walks on G and H starting on vertices g and h , respectively.

Proof The adjacency matrix of $G \oplus H$ is given by $I \otimes H + G \otimes I$. Since $I \otimes H$ and $G \otimes I$ commute, we have

$$|\psi_{G \oplus H}(t)\rangle = e^{-it(I \otimes H + G \otimes I)} |(g, h)\rangle_{G \oplus H} \quad (26)$$

$$= e^{-it(I \otimes H)} e^{-it(G \otimes I)} (|g\rangle_G \otimes |h\rangle_H) \quad (27)$$

$$= (I \otimes e^{-itH}) (e^{-itG} \otimes I) (|g\rangle_G \otimes |h\rangle_H) \quad (28)$$

$$= (e^{-itG} |g\rangle_G) \otimes (e^{-itH} |h\rangle_H) \quad (29)$$

$$= |\psi_G(t)\rangle \otimes |\psi_H(t)\rangle. \quad (30)$$

□

Here, it is convenient to look at the quantum walk on unnormalized adjacency matrices. This is permissible, since normalizing merely affects the time scaling. If k_G and k_H are the regularities of G and H , respectively, Fact 2 may be written as

$$|\psi_{G \oplus H}((k_G + k_H)t)\rangle = |\psi_G(k_G t)\rangle \otimes |\psi_H(k_H t)\rangle \quad (31)$$

in the normalized case. Similar adjustments can be made throughout the rest of this paper.

In quantum mechanics, the natural way to combine two systems is through the tensor product. The previous Fact suggests that the Cartesian graph product serves a similar role for quantum walks.

The following Corollary is the key ingredient to proving our result about Hamming graphs. One direction of the Corollary was already proved in [5].

Corollary 3 *A quantum walk on $G \oplus H$ starting on vertex (g_0, h_0) is instantaneous uniform mixing at time t^* if and only if quantum walks on G and H starting on vertices g_0 and h_0 , respectively, are instantaneous uniform mixing at time t^* .*

Proof Fact 2 shows that the amplitude at vertex (g, h) is given by

$$(\langle g|_G \otimes \langle h|_H)(|\psi_G(t)\rangle \otimes |\psi_H(t)\rangle) = \langle g|\psi_G(t)\rangle \langle h|\psi_H(t)\rangle. \quad (32)$$

The probability at vertex (g, h) is therefore simply

$$P_{G \oplus H}^{(g,h)}(t) = P_G^g(t) P_H^h(t). \quad (33)$$

If P_G and P_H are uniform at time t^* , then clearly so is $P_{G \oplus H}$. If one of P_G or P_H is not uniform at time t^* , then neither is $P_{G \oplus H}$. For suppose, without loss of generality, that $P_G^{g_1} \neq P_G^{g_2}$. Then for any nonzero P_H^h , we have

$$P_{G \oplus H}^{(g_1,h)} = P_G^{g_1} P_H^h \neq P_G^{g_2} P_H^h = P_{G \oplus H}^{(g_2,h)}, \quad (34)$$

showing $P_{G \oplus H}$ is indeed not uniform. \square

With the help of Corollary 3, we may now easily prove our main Theorem of this section.

Theorem 4 *For all $n \geq 1$, a continuous-time quantum walk on the Hamming graph $H(n, q)$ is not instantaneous uniform mixing, unless $q \leq 4$.*

Proof It is known that K_q is not uniform mixing unless $q \leq 4$ [1]. The proof therefore follows immediately from noting that

$$H(n, q) = \underbrace{K_q \oplus \cdots \oplus K_q}_{n \text{ times}}, \quad (35)$$

and repeatedly applying Corollary 3. \square

5 Bunkbed Variants

In this section, we shall view the hypercube Q_{n+1} as a bunkbed graph with adjacency matrix $I \otimes Q_n + X \otimes I$. To this end, we consider a bunkbed operator $\mathcal{B}_n(A)$ defined by

$$\mathcal{B}_n(A) = I \otimes Q_n + X \otimes A = \begin{bmatrix} Q_n & A \\ A & Q_n \end{bmatrix} \quad (36)$$

where A specifies the *connection* between the two copies of Q_n . We investigate the effect of this connection matrix A on the quantum uniform mixing of $\mathcal{B}_n(A)$.

Let the connection graph A be a \mathbb{Z}_2^n -circulant defined by a Boolean function $f : \mathbb{Z}_2^n \rightarrow \{0, 1\}$. The eigenvalues of A are given by $\lambda(a) = \sum_x f(x) \chi_a(x) = \hat{f}(a)$, for $a \in \mathbb{Z}_2^n$. For $a \in \mathbb{Z}_2^{n+1}$, we write $a = a_1 \cdot \tilde{a}$, where $a_1 \in \mathbb{Z}_2$ is the first bit of a and $\tilde{a} \in \mathbb{Z}_2^n$ consists of the remaining bits of a . The eigenvalue λ_a of $\mathcal{B}_n(A)$ is given by

$$\lambda(a) = \lambda_{Q_n}(\tilde{a}) + (-1)^{a_1} \hat{f}(\tilde{a}), \quad (37)$$

In what follows, let $\Delta = \lambda(a \oplus b) - \lambda(b)$ and $\tilde{\Delta} = \lambda(\tilde{a} \oplus \tilde{b}) - \lambda(\tilde{b})$. Using similar analysis to Equation (6), we obtain

$$\hat{P}_t(a) = \frac{1}{2^{n+1}} \sum_b \exp \left[-it \left(\tilde{\Delta} + (-1)^{b_1} \left\{ (-1)^{a_1} \hat{f}(\tilde{a} \oplus \tilde{b}) - \hat{f}(\tilde{b}) \right\} \right) \right] \quad (38)$$

First, we show a simple proof for the hypercube using this bunkbed framework.

Fact 5 (Moore-Russell [11]) *For $n \geq 1$, $Q_{n+1} = \mathcal{B}_n(I)$ is uniform mixing.*

Proof Since $A = I$, we have $f(x) = \llbracket x = 0_n \rrbracket$; thus, $\hat{f}(a) = 1$, for all a . Using this in Equation (38), we get

$$\hat{P}_t(a) = \frac{1}{2^{n+1}} \sum_b \exp \left[-it \left(\tilde{\Delta} + (-1)^{b_1} ((-1)^{a_1} - 1) \right) \right]. \quad (39)$$

Therefore,

$$\hat{P}_t(0 \cdot \tilde{a}) = \frac{1}{2^n} \sum_{\tilde{b}} \exp(-it\tilde{\Delta}) = \hat{P}_t(\tilde{a}) \quad (40)$$

$$\hat{P}_t(1 \cdot \tilde{a}) = \frac{1}{2^{n+1}} \sum_b \exp \left[-it(\tilde{\Delta} - 2(-1)^{b_1}) \right] = \cos(2t) \hat{P}_t(\tilde{a}). \quad (41)$$

By induction, this shows that $\hat{P}_t(a) = \cos(2t)^{|a|}$. Thus, at $t^* \equiv \pi/4$, we have $\hat{P}_{t^*}(a) = \llbracket a = 0_n \rrbracket$, which proves that P_{t^*} is the uniform distribution. \square

In the bunkbed framework, Q_n^η is given by $\mathcal{B}_n(A_f)$, where $\text{supp}(f) = \{0_n, \tilde{\eta}\}$; here, we assume without loss of generality that $\eta = 1 \cdot \tilde{\eta}$. Next, we show that the bunkbed framework is a natural setting for showing limits on uniform mixing.

Theorem 6 *For $n \geq 1$, let $G = \mathcal{B}_n(A_f)$ be a bunkbed hypercube, where A_f is a \mathbb{Z}_2^n -circulant defined by a Boolean function $f : \mathbb{Z}_2^n \rightarrow \{0, 1\}$. Then, $\mathcal{B}_n(A_f)$ is not uniform mixing whenever $|\text{supp}(\hat{f})| < 2^{n-1}$.*

Proof Note that

$$\lambda(a \oplus b) - \lambda(b) = \left[\lambda_{Q_n}(\tilde{a} \oplus \tilde{b}) - \lambda_{Q_n}(\tilde{b}) \right] + (-1)^{b_1} \left[(-1)^{a_1} \hat{f}(\tilde{a} \oplus \tilde{b}) - \hat{f}(\tilde{b}) \right] \quad (42)$$

$$= \left[\lambda_{Q_n}(\tilde{a} \oplus \tilde{b}) - \lambda_{Q_n}(\tilde{b}) \right] - (-1)^{b_1} \left[\hat{f}(\tilde{a} \oplus \tilde{b}) + \hat{f}(\tilde{b}) \right], \quad \text{if } a_1 = 1 \quad (43)$$

Thus, at $a = 1 \cdot 0_n$, we obtain

$$\hat{P}_t(1 \cdot 0_n) = \frac{1}{2^{n+1}} \sum_b \exp \left[2it(-1)^{b_1} \hat{f}(\tilde{b}) \right] = \frac{1}{2^n} \sum_{\tilde{b}} \cos(2\hat{f}(\tilde{b})t) \quad (44)$$

$$= \frac{1}{2^n} \left[2^n - \left(|\text{supp}(\hat{f})| - \sum_{\tilde{b} \in \text{supp}(\hat{f})} \cos(2\hat{f}(\tilde{b})t) \right) \right]. \quad (45)$$

When $|\text{supp}(\hat{f})| < 2^{n-1}$, this is strictly positive, and so P_t cannot be uniform. \square

Remark 3 The condition $|\text{supp}(\hat{f})| < 2^{n-1}$ in Theorem 6 is tight. This is because for $Q_n^\eta = \mathcal{B}_n(A_f)$, where $\text{supp}(f) = \{0_n, \tilde{\eta}\}$ with $\eta = 1 \cdot \tilde{\eta}$, we have $\hat{f}(a) = 1 + \chi_a(\tilde{\eta})$. This shows that $|\text{supp}(\hat{f})| = 2^{n-1}$ and, by Theorem 1, Q_n^η is uniform mixing.

Corollary 7 For $n \geq 2$, a continuous-time quantum walk on the join of two hypercubes $\mathcal{B}_n(J) = Q_n + Q_n$ is not instantaneous uniform mixing.

Proof The connection matrix J is defined by the constant function $f \equiv 1$. Thus, $\hat{f}(a) = \mathbb{I}[a = 0_n]$ and $|\text{supp}(\hat{f})| = 1$. By Theorem 6, $\mathcal{B}_n(J)$ is not uniform mixing. \square

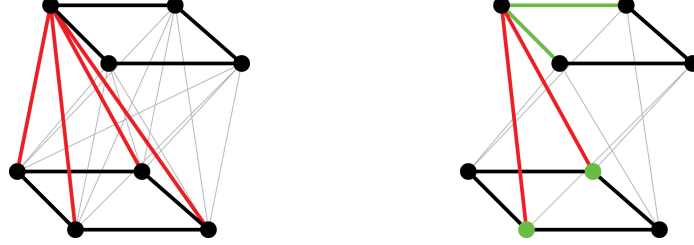


Figure 5: Examples of hypercube bunkbeds: (a) $\mathcal{B}_2(J) = Q_2 + Q_2$. (b) $\mathcal{B}_2(Q_2)$.

The next result shows that $|\text{supp}(\hat{f})| < 2^{n-1}$ in Theorem 6 is not a necessary condition for $\mathcal{B}_n(A_f)$ to be non-uniform mixing.

Theorem 8 For $n \geq 2$, a continuous-time quantum walk on $\mathcal{B}_n(Q_n)$ is not instantaneous uniform mixing.

Proof The adjacency matrix A of $\mathcal{B}_n(Q_n)$ is $J_2 \otimes Q_n$. The eigenvectors of A are $|\chi_a\rangle = |\chi_{a_1}\rangle \otimes |\chi_{\tilde{a}}\rangle$, with corresponding eigenvalues $\lambda(a_1 \cdot \tilde{a}) = (1 + (-1)^{a_1})\lambda_{Q_n}(\tilde{a})$. The amplitude of the quantum walk is given by $\psi_t(a) = \exp(-it(1 + (-1)^{a_1})\lambda_{Q_n}(\tilde{a}))$. Using the Convolution Theorem, we obtain for $a = a_1 \cdot \tilde{a}$ and $b = b_1 \cdot \tilde{b}$:

$$\hat{P}_t(a) = \frac{1}{2^{n+1}} \sum_b \exp \left[-it \left\{ (1 + (-1)^{a_1+b_1})\lambda_{Q_n}(\tilde{a} \oplus \tilde{b}) - (1 + (-1)^{b_1})\lambda_{Q_n}(\tilde{b}) \right\} \right] \quad (46)$$

When $a_1 = 0$, this yields

$$\hat{P}_t(a) = \frac{1}{2} + \frac{1}{2^{n+1}} \sum_{\tilde{b}} \exp \left[-2it \left(\lambda_{Q_n}(\tilde{a} \oplus \tilde{b}) - \lambda_{Q_n}(\tilde{b}) \right) \right] \quad (47)$$

$$= \frac{1}{2} + \frac{1}{2} \cos(2t)^{|\tilde{a}|}. \quad (48)$$

For even $|\tilde{a}|$, $\hat{P}_t(a) \neq 0$. Thus, P_t never equals the uniform distribution. \square

Remark 4 For $\mathcal{B}_n(Q_n)$, we have $f(x) = \mathbb{I}[|x| = 1]$ and hence $\hat{f}(a) = n - 2|a|$. Thus, $|\{a : \hat{f}(a) = 0\}| = \binom{n}{n/2} \sim 2^n / \sqrt{n}$ which implies $|\text{supp}(\hat{f})| \geq 2^{n-1}$. This shows that the condition $|\text{supp}(\hat{f})| < 2^{n-1}$ in Theorem 6 is not a necessary condition.

6 Conclusion

In this paper, we have studied quantum uniform mixing on generalized hypercubes. Our first generalized hypercube is the bunkbed graph $\mathcal{B}_n(A_f)$, where A_f is a \mathbb{Z}_2^n -circulant defined by a Boolean function $f : \mathbb{Z}_2^n \rightarrow \{0, 1\}$. The second generalized hypercube we consider is the Hamming graph $H(n, q)$ or q -ary hypercube. Following the work of Moore and Russell on the hypercube, our main goal is to characterize which of these generalized hypercubes are uniform mixing and which are not.

Prior to this work, there was only one known collection of graphs that is instantaneous uniform mixing. This is the class of Hamming graphs $H(n, q)$ with $q \in \{2, 3, 4\}$ (see [11, 5]). To this collection, we have added $\mathcal{B}_n(A_f)$, where $\text{supp}(f) = \{0_n, a\}$ with $|a|$ odd. This shows that $\mathcal{B}_n(A_f)$ is uniform mixing if $0 < |\text{supp}(f)| \leq 2$ (with the aforementioned condition when $|\text{supp}(f)| = 2$); thus generalizing the result of Moore and Russell [11]. We also showed that if $|\text{supp}(\hat{f})| < 2^{n-1}$, then $\mathcal{B}_n(A_f)$ is not uniform mixing. On the Hamming graphs, we have shown a tight characterization of quantum uniform mixing. Our main result states that $H(n, q)$ is not uniform mixing if and only if $q \geq 5$. This closes the gap left open from an earlier observation in [5].

Unlike the Hamming graphs, the bunkbed framework does not yield a tight characterization, since a counterexample is supplied by $\mathcal{B}_n(Q_n)$. The Fourier techniques we employed do not seem powerful enough to prove a tight characterization based solely on the support sizes of f and \hat{f} . Also, we suspect that quantum uniform mixing on the hypercube Q_n is robust under the addition of up to $O(n)$ specific matchings. We leave these as open problems for future work.

Acknowledgments

This research was supported in part by the National Science Foundation grant DMS-0646847 and also by the National Security Agency grant 42642.

References

- [1] Amir Ahmadi, Ryan Belk, Christino Tamon, and Carolyn Wendler, "On Mixing of Continuous-Time Quantum Walks on Some Circulant Graphs," *Quantum Information and Computation* **3**(6):611-618, 2003.
- [2] William Adamczak, Kevin Andrew, Leon Bergen, Dillon Ethier, Peter Hernberg, Jennifer Lin, and Christino Tamon, "Non-uniform mixing of quantum walk on cycles," *International Journal of Quantum Information* **5**(6):781-793, 2007.
- [3] Norman Biggs, *Algebraic Graph Theory*, 2nd edition, Cambridge University Press, 1993.
- [4] Andrew M. Childs, Richard E. Cleve, Enrico Deotto, Edward Farhi, Sam Gutmann, and Daniel A. Spielman, "Exponential algorithmic speedup by a quantum walk," Proc. 35th Annual Symposium on the Theory of Computing (2003), 59-68.

- [5] William Carlson, Allison Ford, Elizabeth Harris, Julian Rosen, Christino Tamon, and Kathleen Wrobel, "Universal Mixing of Quantum Walk on Graphs," *Quantum Information and Computation* **7**(8):738-751, 2007.
- [6] Persi Diaconis, *Group Representations in Probability and Statistics*, Institute of Mathematical Statistics, 1988.
- [7] Edward Farhi and Sam Gutmann, "Quantum computation and decision trees," *Physical Review A* **58** (1998), 915-928.
- [8] Edward Farhi, Jeffrey Goldstone, and Sam Gutmann, "A Quantum Algorithm for the Hamiltonian NAND Tree," quant-ph/0702144.
- [9] Heath Gerhardt and John Watrous, "Continuous-time quantum walks on the symmetric group," *Proceedings of the 7th International Workshop on Randomization and Approximation Techniques in Computer Science*, 2003.
- [10] Viv Kendon, "Quantum walks on general graphs," *International Journal of Quantum Information* **4**:5 (2006), 791-805.
- [11] Cristopher Moore and Alexander Russell, "Quantum Walks on the Hypercube," in Proc. 6th Int. Workshop on Randomization and Approximation in Computer Science, Lecture Notes in Computer Science **2483**, Springer (2002), 164-178.